Effective and Secure Scheme for Video Streaming Using SRTP

P. Iyyanar, M. Chitra, and P. Sabarinath

Abstract-Real-time of live video or stored video is the predominant part of the real-time multimedia networking. In the streaming technology, the video file need not be downloaded in full, but is being played out while content of the video file are being received and decoded. To send video across the Internet has two approaches: video data compression and design of communication protocol. One approach is to design a low bit rate coder as well as protecting the result of the bit stream using UDP. Another approach is to design new transport protocol by using video coding algorithm. Security plays a major role in streaming application. In this paper, we present an efficient solution for encryption on the video streams using AES-CFB on un-trusted client-server networks. The valuation of the prototype system is to provide a safer transmission across the Internet with cost-effective and quality-guaranteed manner while using SRTP.

Index Terms—Video streaming, AES, RTP, RTSP, SRTP, SRTCP, UDP, H.264.

I. INTRODUCTION

Streaming visual data to different users is becoming ever more popular in recent times, and protecting the transmitted data from every possible security threat has become one of the main concerns both for the end users and data providers. This paper describes a method for protecting streamed data from possible security attacks and suggests a design of secured system architecture for multimedia video streaming to one receiver at a time considering the state of the art for the video streaming existing today. The main feature of the suggested design is its ability to provide a secure communication environment for real-time data.

II. VIDEO STREAMING

Streaming is the process of playing the audio and video file still it downloading. Video streaming [1] refers to the real time transmission of stored video or live video. There are two types for transmission of stored video across the Internet available. One is download mode and another one is streaming mode.

In the download mode, a user downloads the entire video file and then plays back the video file. In the steaming mode, the video file need not be downloaded in full, but is being

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played out while parts of the video file being received.

In real time nature, video streaming has bandwidth, delay and loss of packet requirements. In video streaming, raw video and audio data are compressed by video compression and audio compression algorithms and saved in storage devices. If the client gives the request, a streaming server retrieves compressed audio and video data from storage devices to that particular client. When start to send the audio and video streams across the network the transport protocol packetize the compressed bit streams and send the audiovideo packets to the Internet. For packets that are successfully delivered to the receiver, they first pass through the transport layers and then are processed by the application layer being decoded at the audio- video decoder.

To improve the streaming quality while transmission of audio and video data, continuous media distribution services and media synchronization are developed in the Internet.

III. VIDEO COMPRESSION TECHNIQUE (H.264)

H.264 [2] is a business standard for video compression, the method of converting digital video into a format that takes up a smaller amount capacity when it is stored or transmitted. Video compression is a vital technology for applications such as digital television, Video, mobile TV, videoconferencing and internet video streaming. Customizing video compression makes it possible for products from different producers (e.g. encoders, decoders and storage media) to inter-operate. An encoder transfers video into a compressed format and a decoder convert's compressed video back into an uncompressed format.

IV. REAL-TIME TRANSPORT PROTOCOL

Real-time transport protocol (RTP) [3] is an IP-based protocol providing support for the transport of real-time data such as video and audio streams. RTP offer end-to-end delivery services for data with real-time characteristics. The services provided by RTP consist of time reconstruction, loss detection, security and content identification.RTP itself on the other hand, does not provide all of the functionality required for the transport of data and, therefore, applications usually run it on top of a transport protocol such as UDP. RTP does not deal with resource reservation and does not assurance quality-of-service (QoS) for real-time services. It requests support from lower layers that actually have control over resources in switches and routers.

RTP has its companion RTP control Protocol (RTCP) [5], which is used to control the flow and quality of data and allow the recipient to send feedback to the sources between

Manuscript received June 26, 2012, revised September 01, 2012. This work will be supported in part of Ph.D work.

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client server systems. RTCP has five types messages those are sender report, receiver report, source description message, bye message and application specific message.

V. SECURE REAL-TIME TRANSPORT PROTOCOL

The Secure Real-time Protocol [4] is a profile of the Real-time Transport Protocol (RTP) offering not only confidentiality, but also message authentication, and replay protection for the RTP traffic as well as RTCP (Real-time Transport Control Protocol).

SRTP offers a structure for encryption and message authentication of RTP and RTCP streams. SRTP can achieve high throughput and low packet expansion.

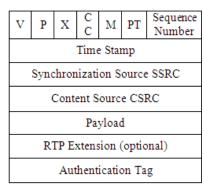


Fig. 1. SRTP Packet format

SRTP is independent of a specific RTP stack implementation and of a specific key management standard, but Multimedia Internet Keying (MIKEY) has been designed to work with SRTP.

In comparison to the security options for RTP there are some advantages to using SRTP. The advantages over the RTP standard security and also over the H.264 security for media stream data are listed below

SRTP provides increased security, achieved by

- Confidentiality for RTP as well as for RTCP by encryption of the respective payloads.
- Integrity for the entire RTP and RTCP packets, together with replay protection.
- The possibility to refresh the session keys periodically, which limits the amount of cipher text produced by a fixed key, variable for an adversary to cryptanalysis.
- An extensible framework that permits upgrading with new cryptographic algorithms.
- A secure session key derivation with a pseudo-random function at both ends.
- The usage of salting keys to protect against precomputation attacks.
- Security for unicast and multicast RTP applications.

VI. PROPOSED SYSTEM FOR SECURED VIDEO STREAMING

In the secured video streaming implementation, the compressed audio-video data is retrieved and packetized at the SRTP layer for the Data plane at the sending side. The SRTP packetized streams provide timing and synchronization information and as well as sequence number. The SRTP packetized streams are then passed to the UDP

layer and the IP layer. The resulting IP packets are transported across the Internet. At the receiver side the media streams are processed in the reversed manner before their presentation. For the control plane, SRTCP packets and RTSP packets are multiplexed at the UDP layer and are moved to the IP layer for transmission across the Internet.

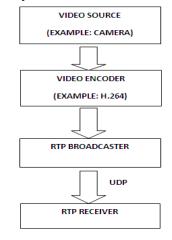


Fig. 2. Video Streaming across the Internet

Fig. 2 shows the Video Streaming across the Internet. In our system we will capture the video from camera or stored audio/ video files and have to encode the video by using H.264 then that video will be separate packet by packet then it will be sent as streams through RTP protocol and that will be received in another side of RTP media player.

A. Architecture of Secured Video Steaming

The big breakthrough that enables the streaming revolution is the adoption of a Internet protocol called the User Datagram Protocol (UDP) and new encoding techniques that compressed audio files into extremely small packets of data. UDP made streaming media possible by transmitting data more powerfully than previous protocols (HTTP and TCP) from the host server over the Internet to the client player or end listener. More recent protocols such as the Real-Time Streaming Protocol (RTSP) [6] are making the transmission of data even more efficient.

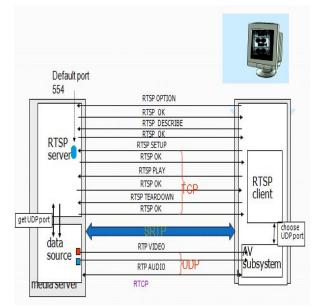


Fig. 3. Video streaming from server

Fig. 3 shows the RTSP client requests the audio and video files from the RTSP server. First we have to implement the RTSP client operation such as RTSP OPTION, RTSP DESCRIPE, RTSP SETUP, RTSP PLAY, RTSP TEARDOWN and RTSP OK for the remote control usage for end-user. Table 1 shows the methods in RTP.

TABLE	I: RTSP METHODS
Methods	Description
Options	Get available methods
ANNOUNCE	Get description of media object
PLAY	Start playback, reposition
RECORD	Start recording
REDIRECT	Redirect client to new server
PAUSE	Halt delivery, but keep state
SET-PARAMETER	Device or encoding control
TEARDOWN	Remove state

In the RTSP and RTP implementation we are using the Wireshark tools for simulation purpose in the client – server connection in which how the packets are received from server to clients. Fig. 5 and Fig. 6 describe the RTSP Client and Server operation.

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Fig. 4. RTSP client connection

The audio and video is compressed and send to the transport layer RTP/UDP then it will go to the IP layer and then transmitted to the Internet. While transmit the audio-video packet to the transport layer the SRTP protocol is used for the packet transaction.

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B. Packet Transmission

There is a number of Internet drafts describe SRTP packetization schemes for MPEG-4 video data. Media aware packetization is a principle in SRTP, so it is likely that several SRTP scheme will be needed to suit the different kind of media, audio, video and so forth. The SRTP time stamp corresponds to the presentation time if the earliest access unit is within the packet. SRTP packets have sequence numbers in transmission order. The payloads logically or physically have synchronization layer sequence numbers, which are decoding order, for each elementary stream.

The MPEG-4 time scale is the time-stamp resolution in the case of MPEG-4 systems and must be used as the SRTP time scale. Streams should be synchronized using RTP techniques which is RTCP sender report. When the MPEG-4 object clock reference is used, it is logically mapped to the network time protocol time axis used in RTCP.

C. AES Chiper

The Rijndael proposed for AES [7] defined a cipher in which the block length and key length can be independently specified to be 128,192 or 256 bits. The AES requirement uses the same three key size alternatives but limits the block length to 128 bits.

Stream cipher techniques have been used in video streaming over the multimedia network. A stream cipher is a symmetric encryption algorithm in which cipher text output is produced bit-by-bit or byte-by-byte from a stream of plaintext input. A block cipher algorithm is used for providing the data security. AES Cipher Feedback (CFB) [8] algorithm is used for general-purpose stream-oriented transmission and provides the authentication. A stream cipher eliminates the need to pad a message to be an integral number of blocks. This is used to operate in real time. In this method a character stream is being transmitted, each character can be encrypted and transmitted immediately using a character-oriented stream cipher.

In the encryption side, the input is a b-bit shift register that is initially set to some initialization vector (IV). The leftmost s bits of the output of the encryption function are XORed with the first segment of plaintext P1 to produce the first unit of cipher text C1, which is then transmitted.

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$$Pi = Ek(Ci-1) \bigoplus Ci \tag{2}$$

$$C0 = IV \tag{3}$$

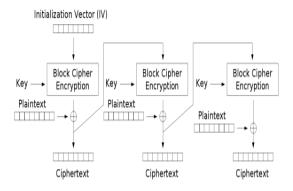


Fig. 6. Cipher Feedback mode encryption

In addition, the contents of the shift register are shifted left by s bits and C1 is placed in the rightmost s bits of the shift register. This process continuous until all the plaintext units has been encrypted.

D. Encryption in SRTP Packet

In our proposed system, we have to modify an encryption method in the SRTP packet format at RTP payload field. In the related work of SRTP implementation, AES – CBC encryption algorithm and AES –CM algorithm were used. In our proposed system, we will apply AES – CFB algorithm for video streams encryption and decryption. The disadvantage of previous algorithm is while encrypting the video file it had encrypted block by block and it gave the general purpose block oriented transmission. In AES –CFB, the stream cipher will be encrypted bit by bit or byte by byte. The advantage of AES – CFB algorithm is to send the stream oriented data transmission and it will give more security than CBC and CM mode in Advanced Encryption Standard.

VII. APPLICATION LAYER QOS

The purpose of application layer QoS control is to avoid congestion and maximize video quality in the presence of packet loss. The application layer QoS control techniques include congestion control and error control. These technique are employed by the end systems and do not require any QoS support from the network. For streaming video, congestion control obtains the form of rate control. There are three kind of rate control: source-based, receiver-based and hybrid rate control.

The source-based rate control is suitable for Unicast video and other two rate control for multicast video.

In unicast video streaming, the model based approach is based on a throughput of a UDP connection. Specifically, the throughput of a UDP connection can be characterized by the following equation:

$$\mu = (1.22 \text{ MTU}) / (\text{RTT } \sqrt{\text{p}})$$
(4)

where μ is a throughput of a UDP connection. MTU is the packet size used by the connection,

RTT is the round-trip time for the connection and p is the packet-loss ratio experienced by the connection. This equation is used to determine the sending rate of the video stream. Thus, the video connection could avoid congestion in a way similar to that of UDP and it can compete fairly with UDP flows.

VIII. THE VALUATION OF SECURE VIDEO STREAMING

In our proposed system, it will provide secure delivery of video data transmission. The time difference between without security of video data transmission and with security of video data transmission will be few seconds. Due to encryption techniques in our system we can find the difference between secure and non secure video streaming transmission. The secured video streaming file will take more time (seconds) than non secure video streaming, but it will provide safer video data transmission in un-trusted client – server networks.

IX. CONCLUSION

RTP server applications transmit captured or stored media streams across the network. The main challenge in designing a video streaming application across the multimedia networks is how to deliver video streams to users with minimal replay jitters with video data security and efficient video data transmission.

The media streams might be encoded in multiple media formats and sent out on several RTP sessions for conferencing with heterogeneous receivers. This paper proposed a framework for video streaming services using SRTP through the client-server network.

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